

# Bridging Quantum Error Correction and Mitigation

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# QEC & QEM

## QEM:

- Low or no qubit overhead.
- Low or no requirement on gate fidelity (few additional quantum operations needed)
- Fast computation on unencoded (or low-distance) qubits

## QEC:

- Exponential suppression of error with increased qubit overhead and without sampling overhead.
- Universal applicability to all algorithms

# Virtual Quantum Error Correction

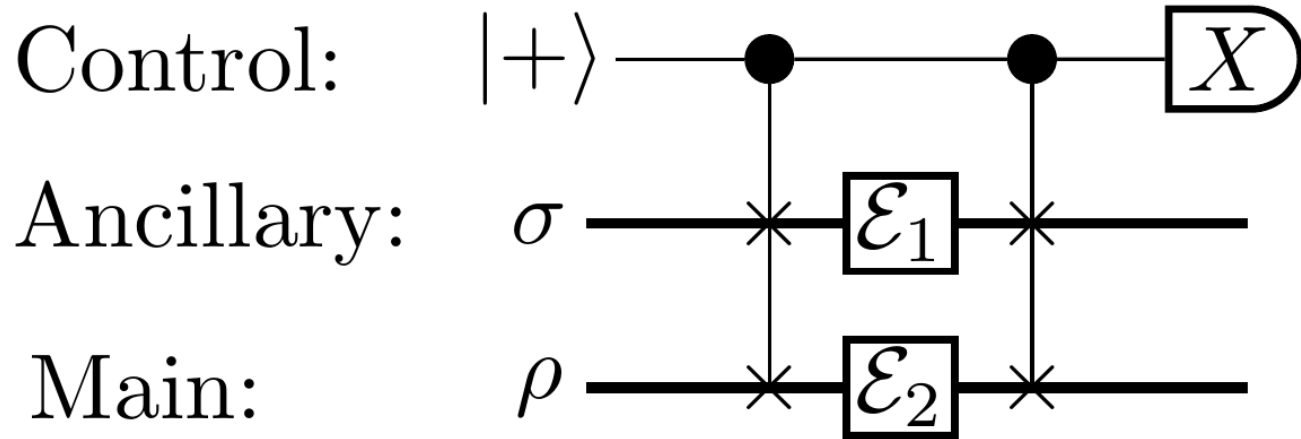
Liu et al, “Virtual Channel Purification”, PRX Quantum 6 (2), 020325

Prior Work:

Piveteau et al, PRL 127, 200505 (2021),

Suzuki et al, PRX Quantum 3, 010345 (2022).

# Virtual Channel Entanglement



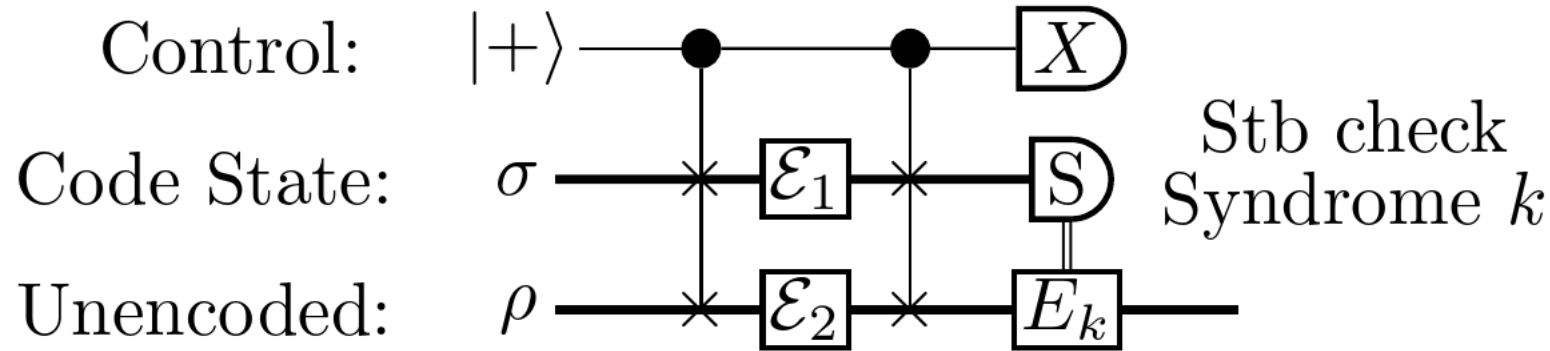
$$\mathcal{E}_{1/2}(\sigma) = \sum_i p_i E_i \sigma E_i^\dagger$$

- After post-processing based on the  $X$  measurement, the effective output state for the two unmeasured registers is

$$\sum_{i,j} p_i^2 (E_j \sigma E_i^\dagger) \otimes (E_i \rho E_j^\dagger)$$

where the two noise channels are virtually entangled.

# Virtual Quantum Error Correction



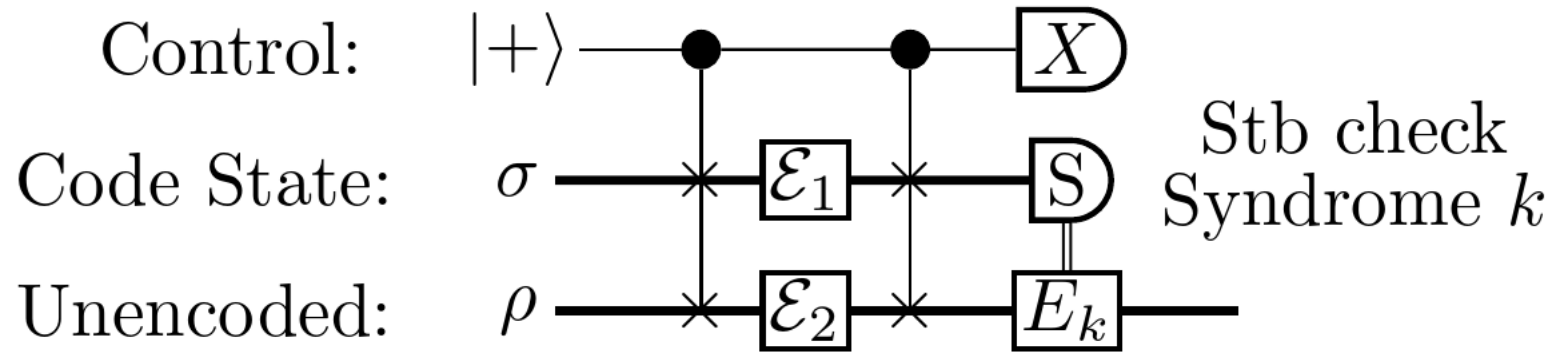
- Choosing  $\sigma$  to be **the code state** of a **non-degenerate** QEC code that **can correct  $\mathcal{E}_{1/2}$**
- Stb measurements  $\rightarrow$  syndrome  $k$ ,  $\rightarrow$  collapse incoming errors into  $E_k$
- We are performing QEC on the **unencoded** register!

$$\sum_{i,j} p_i^2 (E_j \sigma E_i^\dagger) \otimes (E_i \rho E_j^\dagger)$$

$\downarrow$   
 $E_k \sigma E_k^\dagger$

$\downarrow$   
 $E_k \rho E_k^\dagger$   
 $\downarrow$  Correction  $E_k$   
 $\rho$

# Virtual Quantum Error Correction

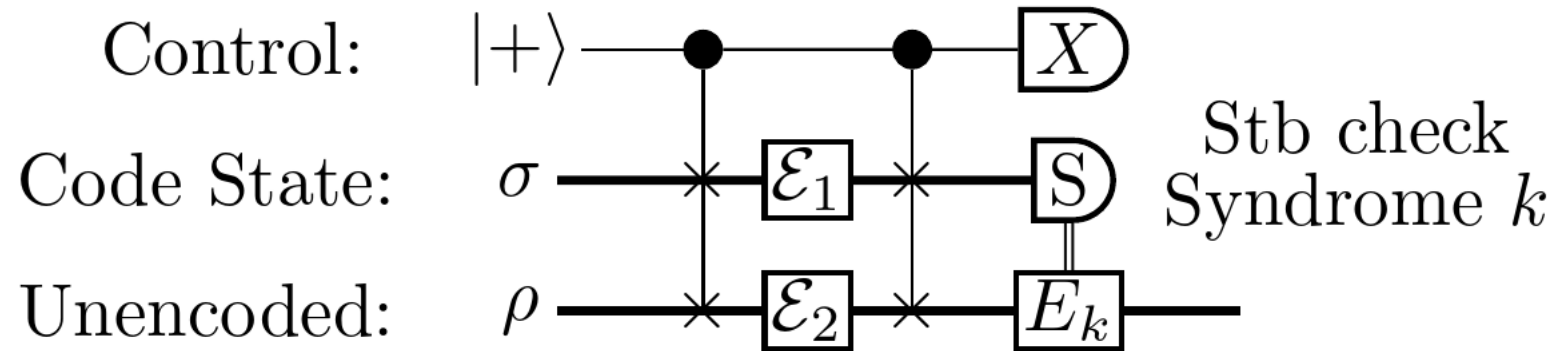


- The error correction is **virtual** because it requires post-processing based on control-qubit  $X$  measurement
- The post-processing comes with a **sampling overhead** of around

$$\sim \left( \sum_k p_k^2 \right)^{-2}$$

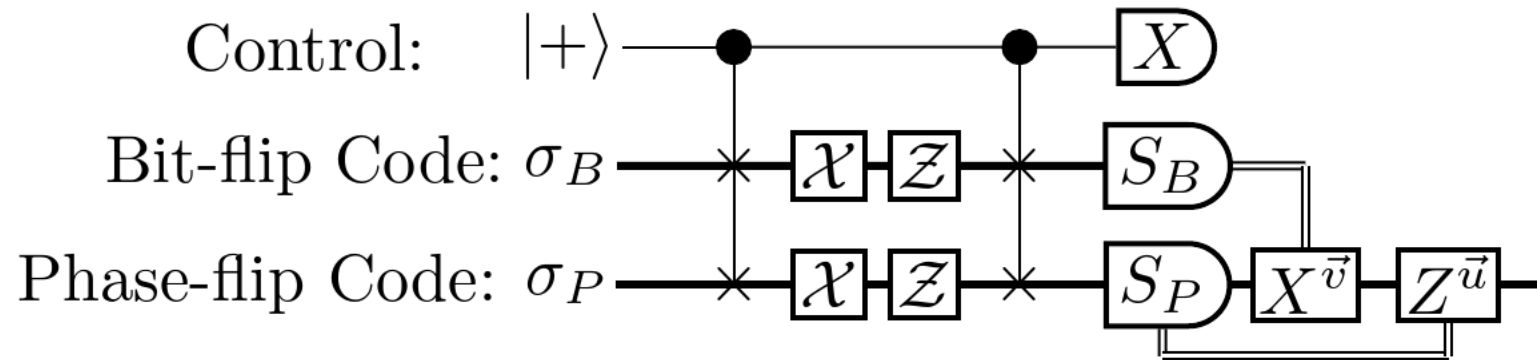
which is similar to virtual state purification (distillation), but now we can achieve **the same noise suppression power as QEC**.

# Comparison to pure QEC



- Goal: send a  $K$ -qubit state  $\rho$  through the noise channel  $\mathcal{E}_2$ .
- Code overhead: suppose we need  $K$  physical qubits per logical qubit
- QEC: the total number of qubit required is  $K^2$ .
- VEC: The total number of qubits required is  $2K + 1$ .
- Limitation: VEC only works when the errors in  $\mathcal{E}_{1/2}$  is **non-degenerate** (distinct syndromes for distinct errors ) for the given code.

# Combining two codes



- Instead of an encoded and an unencoded register, we put one register in bit-flip code and the other in phase-flip code
- Bit-flip check collapses the bit-flip noise on **both registers**, similarly for phase noise.
- So we only need two classical codes to correct quantum noise (both bit-flip and phase-flip)



# Comparison to surface code

- Using **distance- $d$**  bit-flip and phase-flip **repetition codes** for VEC, we can correct the same set of errors as **distance- $d$  surface codes**.
- Advantages:
  - Qubit overhead:  $2d + 1$  (VEC) VS  $d^2$  (surface code)
- Limitation:
  - Sampling overhead: same as virtual channel purification (two-copies)
  - Assuming noiseless CSWAP and stb checks. (Can be further mitigated using PEC)
  - Further work needed for computation.

# QEM for Sampling Algorithm

Liu & Cai, “Quantum Error Mitigation for Sampling Algorithms”,  
arXiv:2502.11285

# QEM for Sampling Algorithm

- QEM use post-processing to combine the output from multiple noisy circuits to obtain the error-mitigated expectation values.
  - The effective damage from noise is only reduced for the entire ensemble of circuit runs.
  - The noise remains unchanged or even increases when zoom individual circuit runs.
- Sampling algorithms (e.g. quantum phase estimation): rely on accurate results for every circuit run, thus seems to be inherently incompatible with QEM (except for those uses post-selection).

# Error-mitigated State

- QEM can also be viewed as trying to extract the error-mitigated “states”  $\rho_{em}$  out of the noisy circuit runs.
- The **error-mitigated** expectation value is given as  $\text{Tr}(O\rho_{em})$ . ( $O$  is the observable of interests)
- The error-mitigated states  $\rho_{em}$  is obtained via **linear combination** of output states from different circuit configurations.
- This covers most mainstream QEM techniques.

# Examples of Error-mitigated States

- Linear Zero-noise extrapolation (can be generalized to Richardson):

$$\rho_p = (1 - p)\rho_0 + p\rho_{err} \Rightarrow \rho_0 = \rho_{em} \propto p_2\rho_{p_1} - p_1\rho_{p_2}$$

- Probabilistic error cancellation for bit-flip noise:

$$\rho = (1 - p)\rho_0 + pX\rho_0X \Rightarrow \rho_0 = \rho_{em} \propto (1 - p)\rho - pX\rho X$$

- Also applicable to other major QEM techniques like virtual purification and symmetry verification.

# State to Distribution is Taking Expectation.

- The output probability of a string  $z$  is the **expectation value** of the projector  $\Pi_z = |z\rangle\langle z|$ .

$$p(z) = \text{Tr}(\Pi_z \rho)$$

- All probabilities of different  $z$  can be **measured simultaneously**.

In a given circuit run, by measuring in the computational basis which output the string  $z'$ , we have obtained one sample for **all**  $\{\Pi_z\}$  with

- one sample of 1 for the  $\Pi_z$  with  $z = z'$
- one sample of 0 for the  $\Pi_z$  with  $z \neq z'$

# QEM for Recovering Output Distribution

- Obtaining error-mitigated distributions  $p_{em}(z) = \text{Tr}(\Pi_z \rho_{em})$  from error-mitigated states  $\rho_{em}$  is efficient.
- Existing mainstream QEM techniques can be used to extract error-mitigated “states”  $\rho_{em}$ , thus can be straightforwardly extended to extract error-mitigated distributions.

# PEC Example

- Probabilistic error cancellation for bit-flip noise:

$$\rho = (1 - p)\rho_0 + pX\rho_0X \Rightarrow \rho_0 = \rho_{em} = \frac{(1 - p)\rho - pX\rho X}{1 - 2p}$$

- Implementation:

1. Sample  $\rho$  and  $X\rho X$  with probability  $(1 - p)$  and  $p$ , respectively.
2. Measure in computation basis  $\{Z_i\}$ , post-process to obtain the set of observables  $\{\Pi_z\}$ .
3. If  $X\rho X$  is sampled, attach minus sign to the output.
4.  $p_{em}(z)$  is estimated by taking the average over all samples of  $\Pi_z$  and renormalise the result with the  $(1 - 2p)^{-1}$  factor.



# Sampling overhead

- Let us consider the trivial observable  $I$ :

$$\hat{I} = \sum_z \hat{\Pi}_z \Rightarrow \text{Var}[\hat{I}] \approx \sum_z \text{Var}[\hat{\Pi}_z] \quad (|\text{Cov}| < 1)$$

- i.e. the variance of estimating a single observable  $I$  is similar to the **total variance** of estimating the probability of all  $z$ , i.e. the entire probability distribution.
- For a given number of circuit runs, the total variance achieved for all entries in the **entire estimated distribution** is actually similar to the variance of **one single observable**.

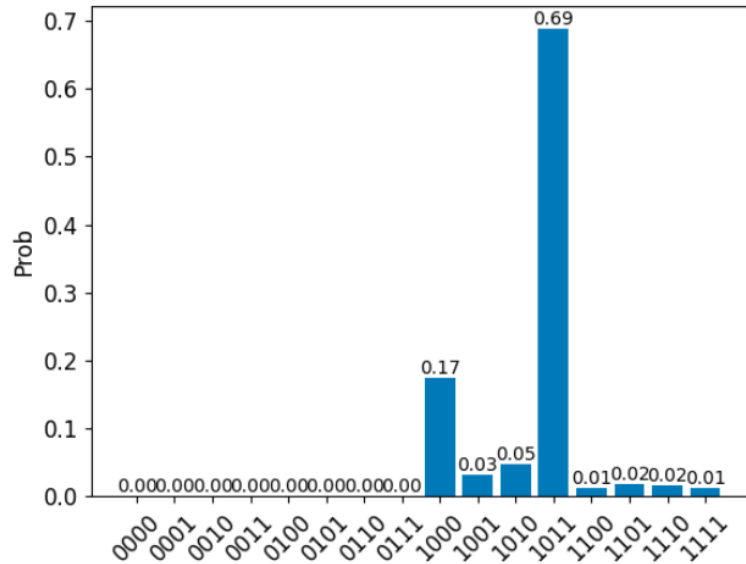
# Application to Quantum Phase Estimation

- Considering using quantum phase estimation for obtaining ground state energy.
- Instead of trying to obtain the whole distribution, we are trying to obtain the smallest string from the output distribution.
- Cannot simply output the smallest string from the estimated error-mitigated distribution, since shot noise can turn zero-probability entries to non-zero.

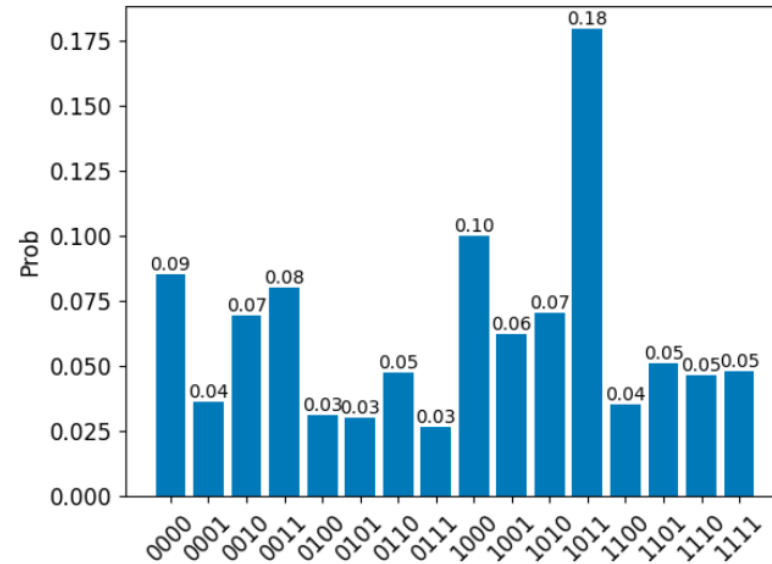
# Application to Quantum Phase Estimation

- An additional step to test whether an entry is likely to be zero or not.
- Set a threshold probability  $p_{\text{th}}(z)$  for each entry such that
  - $\hat{p}_{\text{em}}(z) \leq p_{\text{th}}(z) \Rightarrow \text{Accept null: } p_{\text{em}}(z) = 0$
  - $\hat{p}_{\text{em}}(z) > p_{\text{th}}(z) \Rightarrow \text{Accept null: } p_{\text{em}}(z) > 0$
- $p_{\text{th}}(z)$  can be set using:
  - Proportional to the sample standard deviation of the  $\hat{p}_{\text{em}}(z)$  estimator.
  - Known lower bound of the probability of the smallest string.

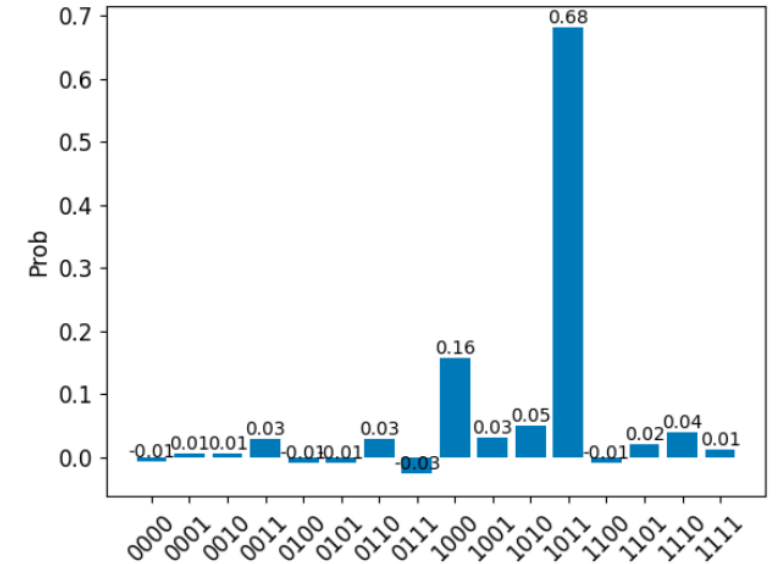
# Numerical Simulation



(a) Ideal  $p_0(z)$



(b) Noisy  $p(z)$



(c) Mitigated  $\hat{p}_{\text{em}}(z)$ ,  $10^6$  circuit runs

- QPE with 4-bit precision
- Circuit error rate  $\sim 0.6$
- $10^6$  runs

- Total square errors reduced from 0.297 to 0.004
- Valid threshold:  $0.03 < p_{\text{th}} < 0.16$

# How to sample from the QEM distribution?

- Without QEM, when measure  $z$  in a circuit run, we put one sample into the “bucket” corresponding to outcome  $z$ .
- With QEM, when measure  $z$  in a circuit run, there is also a additional sign associated with the circuit configuration we are running:
  - +ve sign: **add** one sample into the “bucket” corresponding to outcome  $z$
  - -ve sign: **remove** one sample from the “bucket” corresponding to outcome  $z$

# How to sample from the QEM distribution?

- There can be negative number of samples! Esp. when the number of circuit run is small.
- When comes to interpretation of results, these negative number can effectively be treated as zero since any components below zero are entirely due to shot noise.

# Summary

- VEC is a framework to combine the power of different QEC codes using post-processing.
- QEM can be used for recovering the output distribution and also sampling from it, and this is as cheap as one observable.

## **Open questions:**

- More general frameworks to combine QEC and QEM, hopefully introducing a range of techniques with different trade-offs.
- What are the other (practical) algorithms that still fall outside the remit of QEM? How to include them?